### Probability and Random Processes Lecture 12

- Detection
- Estimation
- Capacity
- Information

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## Detection

- A random process on  $(E^T, \mathcal{E}^T)$  with two possible distributions  $\mu_0$  and  $\mu_1$
- Assume that  $\mu_0 \gg \mu_1$  and  $\mu_1 \gg \mu_0$  (the distributions are equivalent)
- Observe  $f \in E^T$  and based on the observation

decide  $H_0: \mu_0$  or  $H_1: \mu_1$ 

 $\iff$  design a measurable mapping  $g: E^T \to \{0, 1\}$ 

#### Criteria

• Classical: minimize

$$P(g(f) = 0|H_1)$$

subject to  $P(g(f) = 1|H_0) \leq \alpha$ 

• Bayesian: minimize

$$P_e = P(g(f) = 0|H_1)P(H_1) + P(g(f) = 1|H_0)P(H_0)$$

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#### Bayesian detection

• Let  $G_1 = g^{-1}(\{1\})$ ,  $G_0 = g^{-1}(\{0\})$ , assuming  $G_1 \cup G_0 = E^T$ and  $G_1 \cap G_0 = \emptyset$ , then

$$P_e = P(H_1) \int_{G_0} d\mu_1 + P(H_0) \int_{G_1} d\mu_0$$
  
=  $P(H_1) \int_{G_0} \left( \frac{d\mu_1}{d\mu_0} - \frac{P(H_0)}{P(H_1)} \right) d\mu_0 + P(H_0)$ 

• Hence, we should set

$$G_0 = \left\{ f : \frac{d\mu_1}{d\mu_0}(f) < \frac{P(H_0)}{P(H_1)} \right\}$$
$$G_1 = \left\{ f : \frac{d\mu_1}{d\mu_0}(f) > \frac{P(H_0)}{P(H_1)} \right\}$$

• Compare the likelihood ratio

$$\lambda(f) = \frac{d\mu_1}{d\mu_0}(f)$$

to a threshold

- Classical  $\Rightarrow$  Neyman–Pearson: also based on comparing  $\lambda$  to a threshold
- Given  $f \in E^T$ , how do we compute  $\lambda(f)$ ?

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#### Grenander's theorem

- Look at  $(E, \mathcal{E}) = (\mathbb{R}, \mathcal{B})$  and  $T = \mathbb{R}^+$
- $X : (\Omega, \mathcal{A}, P) \to (\mathbb{R}^T, \mathcal{B}^T, \mu)$  is separable if there is a set  $N \subset \Omega$  for which P(N) = 0 and a sequence  $S = \{t_k\} \subset T$  such that for any open interval I and closed set C

$$\{\omega: \pi_t(X(\omega)) \in C, \ t \in I \cap T\} \setminus \{\omega: \pi_t(X(\omega)) \in C, \ t \in I \cap S\} \subset N$$

- i.e.,  $f(t) \in \mathbb{R}^T$  can be sampled without loss at the  $t_k$ 's
- For any  $X : (\Omega, \mathcal{A}, P) \to (\mathbb{R}^T, \mathcal{B}^T, \mu)$  there is a  $\tilde{X} : (\Omega, \mathcal{A}, P) \to (\mathbb{R}^T, \mathcal{B}^T, \tilde{\mu})$  such that  $\tilde{X}$  is separable and

$$P(\{\omega: \pi_t(X) = \pi_t(\tilde{X}), t \in T\}) = 1$$

- Consider the detection problem, and assume that  $\mu_1$  and  $\mu_2$ when restricted to  $\{t_k\}_1^n$ , for any finite n, are both absolutely continuous with respect to Lebesgue measure on  $\mathbb{R}^n$
- Observe f(t), sample as  $f_k = f(t_k)$  (with  $\{t_k\}$  as in the definition of separability), let  $f^n = (f_1, \ldots, f_n)$  and denote the densities  $g_1(f^n)$  and  $g_2(f^n)$  (corresponding to  $\mu_1$  and  $\mu_2$ )
- Then the entity

$$g_n = \frac{g_1(f^n)}{g_2(f^n)}$$

converges with probability one to  $\lambda(f)$  under both  $H_0$  and  $H_1$ 

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Gaussian waveforms

- Consider the continuous-time Gaussian example with two possible mean-value functions, that is, f(t) is Gaussian with  $E[f(t)] = m_i(t)$  under  $H_i$ , and has a positive-definite covariance kernel V(s, u) (under both  $H_0$  and  $H_1$ )
- Without loss we can assume  $m_0(t) = 0$  and  $m_1(t) = m(t)$
- Assume that m(t) can be expressed as

$$m(t) = \int V(t,s)h(s)ds$$

for some h(t), then

$$\ln \lambda(f) = \int f(t)h(t)dt - \frac{1}{2}\int m(t)h(t)dt$$

(with probability one under  $H_0$  and  $H_1$ )

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## Estimation

#### Bayesian

- Two random objects X and Y on (Ω, A, P) with range spaces (ℝ, B) and (E, E) (standard)
- Estimate  $X \in \mathbb{R}$  from observing  $Y = y \in E$ ; MMSE  $\Rightarrow$

$$\hat{X}(y) = E[X|Y = y]$$

• That is,

$$\hat{X}(y) = \int x d\mu_y$$

where  $\mu_y$  is the regular conditional distribution for X given Y=y

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### Classical

• For an absolutely continuous random variable X with pdf  $f_{\alpha}(x)$ ; given the observation X = x we have the traditional ML estimate

$$\hat{\alpha} = \arg\max_{\alpha} f_{\alpha}(x)$$

- A pdf f(x) is the Radon–Nikodym derivative of the distribution μ on (ℝ, ℬ) w.r.t. Lebesgue measure λ; that is, it can be interpreted as the likelihood ratio between the hypothesis H<sub>0</sub> : μ = λ and H<sub>1</sub> : μ = μ<sub>X</sub> (the correct distribution)
- In the case of a general random object X : (Ω, A, P) → (E, E, μ), we can choose a "dummy hypothesis" H<sub>0</sub> : μ = μ<sub>0</sub> as a reference to H<sub>1</sub> : μ = μ<sub>α</sub>, where μ<sub>α</sub> is the correct distribution, with an unknown parameter α ∈ ℝ

• Then, based on the observation X = x, the ML estimate can be computed as

$$\hat{\alpha} = \arg\max_{\alpha} \frac{d\mu_{\alpha}}{d\mu_0}(x)$$

- The reference distribution can be chosen e.g. such that computing the likelihood ratio is feasible
- Note that in general, the estimate "
   *α* = arg max μ<sub>α</sub>" does not make sense; μ<sub>α</sub> is a mapping from sets in *E*

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# Channels

- Given two measurable spaces (Ω, A) and (Γ, S), a mapping g : Ω × S → ℝ<sup>+</sup> is called a transition kernel (from Ω to Γ) if
  - **1**  $f(\omega) = g(\omega, S)$  is measurable for any fixed  $S \in S$ **2**  $h(S) = g(\omega, S)$  is a measure on  $(\Gamma, S)$  for any fixed  $\omega \in \Omega$

If the measure h(S) in 2. is a probability measure, then g is called a stochastic kernel

 If Y is a random object on (Ω, A, P) with values in (E, E), then a stochastic kernel g from Ω to E is the regular conditional distribution of Y given G ⊂ A if

$$g(\omega, F) = P(\{Y \in F\} | \mathcal{G})(\omega)$$

with probability one w.r.t. P and  $\omega,$  and for all  $F\in\mathcal{E}$ 

• A regular conditional distribution for Y exists if  $(E, \mathcal{E})$  is standard

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- The factorization lemma: Assume two measurable spaces

   (Ω<sub>1</sub>, A<sub>1</sub>) and (Ω<sub>2</sub>, A<sub>2</sub>) and a measurable mapping
   u: Ω<sub>1</sub> → Ω<sub>2</sub> are given. A function v : (Ω<sub>1</sub>, A<sub>1</sub>) → (ℝ, B) is
   measurable w.r.t σ(u) ⊂ A<sub>1</sub> iff there is a measurable mapping
   φ: (Ω<sub>2</sub>, A<sub>2</sub>) → (ℝ, B) such that v = φ ∘ u
- Let X be an arbitrary random object on  $(\Omega, \mathcal{A}, P)$ , and let Y be as before, with  $(E, \mathcal{E})$  standard. Let

$$g(\omega, F) = P(\{Y \in F\} | \sigma(X))(\omega)$$

and let  $\phi_F$  be the mapping in the factorization lemma  $g(\omega, F) = \phi_F(X(\omega))$  (w.r.t.  $\omega$  for a fixed F), then

$$P(\{Y \in F\} | X = x) = \phi_F(x)$$

is the conditional distribution of Y given X = x

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- Assume  $(\Omega, \mathcal{A}, P)$ , a parameter set T and a process  $X : (\Omega, \mathcal{A}, P) \to (E^T, \mathcal{E}^T, \mu_X)$  are given
- Given another function/sequence space  $(F^T, \mathcal{F}^T)$ , a channel is a regular conditional distribution from  $\Omega$  to  $F^T$  given a specific value  $X = x \in E^T$ ; that is, for any  $x \in E^T$  the distribution

$$\mu_x(F) = P(\{Y \in F\} | X = x), \ F \in \mathcal{F}^T$$

- Interpretation: A random channel input X is generated and is then transmitted over the channel, resulting in the channel output Y; given X = x the distribution for Y is μ<sub>x</sub>
- A channel exists if the relevant spaces are standard

- Given  $(E^T, \mathcal{E}^T)$  and  $(F^T, \mathcal{F}^T)$ , let  $\mathcal{E}^T \times \mathcal{F}^T$  be the product  $\sigma$ -algebra on  $E^T \times F^T$
- Let  $\tilde{\mu}$  be defined by

$$\tilde{\mu}(A,B) = \int_A P(B|X=x) d\mu_X(x)$$

on rectangles,  $A \in \mathcal{E}^T, B \in \mathcal{F}^T$ 

- Standard spaces  $\Rightarrow$  unique extension of  $\tilde{\mu}$  from rectangles to  $\mathcal{E}^T \times \mathcal{F}^T$ ; a joint distribution  $\mu$  on  $(E^T \times F^T, \mathcal{E}^T \times \mathcal{F}^T)$
- Also define the corresponding product distribution  $\pi$ , generated as the extension of  $\mu_X(A)\mu_Y(B)$ , with

$$\mu_Y(B) = \int_{\Omega} P(B|X=x) d\mu_X(x)$$

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Channel Capacity

- Focus on  $T = \mathbb{N}^+$  and  $E = F = \mathbb{R}$ ; a channel  $\mu_x(\cdot)$  with input  $x \in \mathbb{R}^T$  and output  $y \in \mathbb{R}^T$  (sequences), resulting in the joint distribution  $\mu$
- A rate R [bits per channel use] is achievable if information can be transmitted at R with error probability below  $\varepsilon$  for any  $\varepsilon > 0$
- The capacity C of the channel  $= \sup\{R : R \text{ is achievable }\}$

• Let  $S_n = \{1, 2, \dots, n\}$ , define the information density

$$i(x,y) = \log \frac{d\mu}{d\pi}$$

(assuming  $\pi \gg \mu$ ), and the corresponding restricted version

$$i_n(x^n, y^n) = \log \frac{d\mu_{|S_n|}}{d\pi_{|S_n|}}$$

• Let

$$\gamma(\mu_X) = \sup\left\{\alpha : \lim_{n \to \infty} P(n^{-1}i_n \le \alpha) = 0\right\}$$

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• A general formula for channel capacity [Verdú-Han, '94]:

$$C = \sup_{\mu_X} \gamma(\mu_X)$$

• Computing C involves the problem of characterizing the limit  $\gamma$  (for each fixed  $\mu_X$ )  $\Rightarrow$  ergodic theory

### Information Measures

- Given  $(\Omega, \mathcal{A})$ , a measurable partition of  $\Omega$  is a finite collection  $G_1, \ldots, G_n$ ,  $G_i \in \mathcal{A}$ , such that  $G_k \cap G_l = \emptyset$  for  $k \neq l$  and  $\cup_i G_i = \Omega$
- Given two probability measures P and Q on  $(\Omega, A)$  and a measurable partition  $\mathcal{G} = \{G_i\}_{i=1}^n$  of  $\Omega$ , define

$$D^*(P||Q)(\mathcal{G}) = \sum_{G \in \mathcal{G}} P(G) \log \frac{P(G)}{Q(G)}$$

• Then, the relative entropy between P and Q is defined as

$$D(P||Q) = \sup_{\mathcal{G}} D^*(P||Q)(\mathcal{G})$$

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• If  $P \ll Q$ , then we get

$$D(P||Q) = \int \log \frac{dP}{dQ}(\omega) \, dP(\omega)$$

- Let X and Y be two random objects on (Ω, A, P) with range spaces (Γ, S) and (Λ, U), and a joint distribution μ<sub>XY</sub> on (Γ × Λ, S × U) corresponding to the marginal distributions μ<sub>X</sub> and μ<sub>Y</sub>
- Let  $\pi_{XY}$  be the corresponding product distribution
- The mutual information between X and Y is then defined as

 $I(X;Y) = \sup_{\mathcal{F}} D^*(\mu_{XY} || \pi_{XY})(\mathcal{F})$ 

over measurable partitions  ${\cal F}$  of  $\Gamma \times \Lambda$ 

• The entropy of the single variable X is defined as

H(X) = I(X;X)

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• If  $\mu_{XY} \ll \pi_{XY}$ , then

$$I(X;Y) = \int \log \frac{d\mu_{XY}}{d\pi_{XY}} d\mu_{XY}$$

• Returning to the setup of transmission over a channel (with the previous notation), if  $\pi \gg \mu$  we had

$$i_n(x^n, y^n) = \log \frac{d\mu_{|S_n|}}{d\pi_{|S_n|}}$$

• If for any fixed stationary and ergodic input, with distribution  $\mu_X$ , the channel is such that the joint input-output process on  $(\mathbb{R}^T \times \mathbb{R}^T, \mathcal{B}^T \times \mathcal{B}^T)$  is stationary and ergodic, and in addition satisfies the finite-gap information property (below), then

$$\frac{1}{n}i_n \to \lim_{n \to \infty} \frac{1}{n}I(X^n;Y^n)$$

with probability one

• finite-gap information: for any n > 0 there is a  $k \ge n$  such that  $I(X_k; X^n | X^k)$  and  $I(Y_k; Y^n | Y^k)$  are both finite

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• Letting

$$i_{\infty}(\mu_X) = \lim_{n \to \infty} \frac{1}{n} I(X^n; Y^n)$$

for any fixed  $\mu_X$ , we get

$$C = \sup_{\mu_X} i_{\infty}(\mu_X)$$

- Channels that result in this formula for C have been called information stable
- To prove this, one first needs to see that  $\gamma = i_{\infty}$  for any fixed  $\mu_X$  such that the input, output and joint input-output are stationary and ergodic. Then one needs to show that the supremum is achieved in this class.