

## Differentially Private State Estimation in Distribution Networks with Smart Meters

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### **Motivation**

- **The promise:** Smart meters help in demand response, billing, etc.
- Few real-time measurements in today's distribution networks → Enabler for state estimation?
- The threat: Customers' privacy concerns (among others)
- The opportunity: Privacy-preserving monitoring and control techniques







## **Distribution Network Model**



• Conservation of currents:

$$I_j = \sum_{k>j} L_k, \quad j = 0, 1, \dots, N-1$$

- Large resistance in distribution grids  $\rightarrow$  use currents
- Little dynamics in (current) distribution grids  $\rightarrow$  study steady-state
- **Operator:** Desires to estimate load *L<sub>j</sub>*
- Customer C: Desires to keep his/her real-time load private



### **Measurement Model**



The Base Scenario – Total current with "physical meter noise":

 $Z_0 = I_0 + W_0 \qquad \qquad W_0 \sim \mathcal{N}(0, R_0)$ 

The Smart Meter Scenario – Load current with "privacy noise":

$$Z_j = L_j + W_j \qquad W_j \sim \operatorname{Lap}(b_j)$$
$$p_{W_j}(w) = \frac{1}{2b_j} e^{-|w|/b_j}, R_j = 2b_j^2$$



## **Problem Formulation: Characterize Estimation vs. Privacy Trade-Offs**





#### **Related Work**

- Differential privacy:
  - Dwork, McSherry, Nissim, Smith, 2006
- Differential privacy in control:
  - Le Ny, Pappas, 2014
  - Huang, Wang, Mitra, Dullerud, 2014
- Privacy for Smart Meters:
  - Ács, Castelluccia, 2011
  - Tan, Gunduz, Poor, 2013



## Differential Privacy [Dwork et al., 2006]

• Two adjacent data vectors:

$$l = \begin{pmatrix} l_1 & l_2 & \dots & l_i & \dots & l_{m-1} & l_m \end{pmatrix}^T$$
$$l' = \begin{pmatrix} l_1 & l_2 & \dots & l_i \pm \Delta & \dots & l_{m-1} & l_m \end{pmatrix}^T$$

• Measurement policy (q deterministic, W stoch. noise) Z(l,W) = q(l) + W(We will use  $q(l) = \sum_k l_k$ )

**Definition**: Measurement *Z* is  $(\epsilon, \delta)$ -differentially private if for all events *E*:

 $\Pr[Z(l, W) \in E] \le e^{\epsilon} \Pr[Z(l', W) \in E] + \delta$ 



# Example: $\epsilon$ -Differential Privacy with Laplacian Noise



 $W \sim \operatorname{Lap}(b)$  $\epsilon = \frac{\Delta}{b} = \frac{\Delta}{\sqrt{R/2}}$ 

Measurements of adjacent data vectors virtually indistinguishable for small  $\epsilon$ 



#### **Differential Privacy in the Distribution** Network



-2

-3

-1

0

0.1

0i -4

**The Base Scenario** ( $Z_0$ ): Customer C has ( $\epsilon_0, \delta_0$ )-differential privacy where



 $\delta_0$ 

2

3

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differential



## **Optimal Estimate: Load Model**



Suppose loads have a known normal distribution:  $L \sim \mathcal{N}(m, P)$ 

$$m = \begin{pmatrix} m_1 \\ \vdots \\ m_N \end{pmatrix} \quad P = \begin{pmatrix} P_{11} & \dots & P_{1N} \\ \vdots & \ddots & \vdots \\ P_{N1} & \dots & P_{NN} \end{pmatrix} \quad \begin{pmatrix} P_1 \\ \vdots \\ P_N \end{pmatrix} \coloneqq \begin{pmatrix} P_{11} + \dots + P_{1N} \\ \vdots \\ P_{N1} + \dots + P_{NN} \end{pmatrix} = P\mathbf{1}$$



## **Optimal Estimate: Base Scenario**



**MMSE** estimate:

$$\hat{L}_{j}^{0} := \mathbf{E}[L_{j}|Z_{0}] = m_{j} + \frac{P_{j}}{P_{0} + R_{0}}(Z_{0} - m_{0})$$

**MMSE error:** 

$$Q_j^0 := \mathbf{E}[(\hat{L}_j^0 - L_j)^2] = P_{jj} - \frac{P_j^2}{P_0 + R_0}$$



## **Optimal Estimate: Smart Meter Scenario**



LMMSE estimate:

$$\hat{L}_{j}^{0,j} := \mathbf{E}^{\text{lin}}[L_{j}|Z_{0}, Z_{j}]$$
$$= \hat{L}_{j}^{0} + K_{j} \left[ (Z_{j} - m_{j}) - \frac{P_{j}}{R_{0} + P_{0}} (Z_{0} - m_{0}) \right]$$

LMMSE el 
$$_{K^{+}} - \frac{(R_0 + P_0)P_{jj} - P_j^2}{Q_j^{0,j} := \mathbf{E}[(L_j - \hat{L}_j^{0,j})^2] = Q_j^0(1 - K_j) \le Q_j^0$$



## Trade-Off: Estimation Quality vs. Privacy

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#### **Dimensionless quantities:**

- Customers' relative importance at site *j*:  $\eta_j := \frac{\Delta^2}{P_{ij}}$ •
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Site *j*'s relative importance on the line:  $\zeta_j := \frac{P_{jj}}{P_0 + R_0}$ 



#### **Trade-Off: Estimation Quality vs. Privacy**



Baseline privacy:  

$$\epsilon_0^2 \approx \frac{\Delta^2 K^2}{R_0}$$

$$= \eta_j \zeta_j K^2 \left( 1 + \frac{P_0}{R_0} \right)$$

Est. improvement:

K

$$\begin{aligned} \zeta_j &= \frac{1}{1 + \frac{2\eta_j}{\epsilon^2(1-\zeta_j)}} \\ &\approx \frac{\epsilon^2(1-\zeta_j)}{2\eta_j} \end{aligned}$$



#### Summary



- Simple analytical treatment of trade-off between state estimation quality and customers' privacy loss  $\epsilon$
- Estimation gain  $\sim \left(\frac{\epsilon}{\epsilon_0}\right)^2 \rightarrow$  Customers with high baseline privacy can • make a large difference!
- Possible extensions: Dynamics, general topologies, active/reactive • power flows